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Methods and problems in discrete mathematics

Wintersemester 2019/20

— Exercise Sheet 2 —

**Exercise 2.1** Make the following statement precise: The linear programming bound for potential energy generalizes the linear programming bound for spherical codes<sup>1</sup>.

$$\begin{array}{ll} \inf & \lambda \\ & \beta_0, \beta_1, \dots, \beta_d \geq 0 \\ & \sum\limits_{k=0}^d \beta_k = \lambda - 1 \\ & \sum\limits_{k=0}^d \beta_k P_n^k(t) \leq -1 \text{ for } t \in [-1, \cos \gamma]. \end{array}$$

*Hint:* If you're stuck, a look into the Cohn-Kumar paper might help.

**Exercise 2.2** Consider the sequence  $(p_n)$  of orthogonal polynomials satisfying the orthogonality relation

$$\int_{-1}^{1} p_m(x) p_n(x) \, dx = 0 \text{ for } m \neq n,$$

normalized by  $p_n(1) = 1$ . Determine  $p_0, p_1, p_2, p_3, p_4$  and verify the 3-term recurrence relation for  $p_2, p_3$  and  $p_4$ .

**Exercise 2.3** Determine H(f,g) for

$$f(t) = 1/(2-2t)^3$$
 and  $g(t) = (t+1)^2(t+1/2)^2t^2(t-1/2)^2$ .

**Exercise 2.4** Present the proof of Proposition 2.2 (see page 110 in the Cohn-Kumar paper): Let f be absolutely monotonic in (a, b), let  $g(t) = (t - t_1)^{k_1} \cdots (t - t_m)^{k_m}$ . Then, Q(f,g)(t) is again absolutely monotonic in (a, b).

"Hand-in": Until Thursday October 24, 10 am, using the form on the course homepage.

<sup>&</sup>lt;sup>1</sup>A spherical code with minimal angle  $\gamma$  is a point configuration  $\mathcal{C} \subseteq S^{n-1}$  so that  $x \cdot y \in [-1, \cos \gamma]$  whenever  $x \neq y$ .