

Methods and problems in discrete mathematics

WS 2019/2020

Goals

- Develop a mathematical toolbox (especially algebraic & analytic) to solve problems in discrete mathematics (especially discrete geometry & graph theory).
 - Explain and discuss (open) problems.
 - Get familiar with contemporary literature in discrete mathematics
 - Don't care^{at all} about "purity of methods"
- (→ cross-disciplinary approach)
- Presentation will not be "linear" but rather "hilly"

Chapter I Universally optimal distribution of points

(following a paper by Ghu & Kumar)

Potential energy minimisation

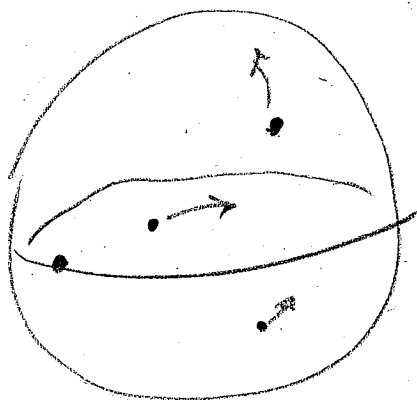
M (compact) Riemannian manifold, for example

$M = S^{n-1} = \{x \in \mathbb{R}^n : \|x\| = 1\}$ the unit sphere.

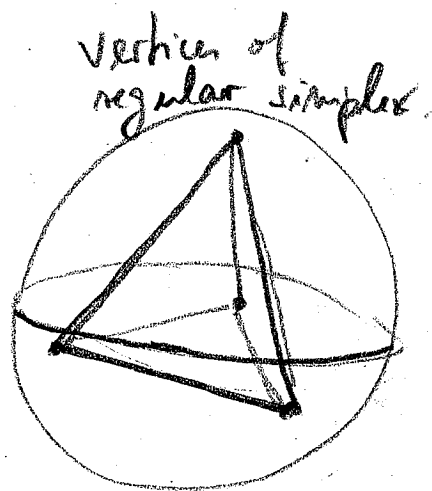
$\mathcal{C} \subseteq M$ finite point set

Question Identically charged particles $\mathcal{C}$ move on M according to a potential function. How do they arrange? physics $\leadsto$ locally minimise their electrical potential energy

S^2



$|\mathcal{C}| = 4$



vertices of regular simplex

Here Focus on global minima.

Def. Let $f: (0, 4] \rightarrow [0, \infty)$ be a decreasing, continuous function (radial potential function).

The f -potential energy of C is

$$E_f(C) = \sum_{\substack{x, y \in C \\ x \neq y}} f(\|x - y\|^2)$$

Example

$$m = 3 \quad f(r) = \frac{1}{r^{1/2}} \quad \text{electrical (Coulomb) potential}$$

$$m \geq 3 \quad f(r) = \frac{1}{r^{m/2-1}} \quad \text{Riesz potential; inverse power law}$$

Natural class of radial potential functions

Def. $I \subseteq \mathbb{R}$ interval

$f: I \rightarrow \mathbb{R}$ C^∞ -function is called

completely monotonic if $(-1)^k f^{(k)}(x) \geq 0$
for all $x \in I, k \in \mathbb{N}$.

$k=0$: $f(x) \geq 0 \rightarrow f$ nonnegative

$k=1$: $-f'(x) \geq 0 \rightarrow f$ decreasing

$k=2$: $f''(x) \geq 0 \rightarrow f$ convex

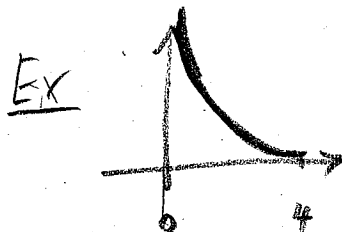
$k=3$: $-f'''(x) \geq 0 \rightarrow f$ has negative jerk
(dt. "Ruck")

$\vdots$

$\leadsto$ causes motion sickness.

Example

(a) inverse power laws



$f(x) = \frac{1}{x^s}$, $s > 0$, are (strictly) completely

monotonic $\uparrow$
all inv. strict
in the interior of I

(b) Gaussian potential

$$f(x) = e^{-cx}, \quad c > 0$$

Theorem Completely monotonic functions form a convex cone. This cone is generated by the functions $x \mapsto (4-x)^k$, $k=0,1,2,\dots$ (if $I=(0,4]$).

Proof see Widder, 1941.

Def. A finite set $\mathcal{C} \subseteq S^{m-1}$ is called a spherical t -design if for every polynomial $p \in \mathbb{R}[x_1, \dots, x_m]$ of degree $\leq t$:

$$\frac{1}{|\mathcal{C}|} \sum_{x \in \mathcal{C}} p(x) = \int_{S^{m-1}} p(x) d\omega(x),$$

where ω is the unique rotationally invariant probability measure on S^{m-1} .

Def. $\mathcal{C} \subseteq S^{m-1}$ is called a sharp configuration if

(a) there are m inner products between distinct pairs of points:

$$|\{x \cdot y : x, y \in \mathcal{C}, x \neq y\}| \geq m$$

(b) $\mathcal{C}$ is a spherical $(2m-1)$ design.

Example vertices of regular simplex form a sharp configuration with $m=1$.

Main Theorem Let $f: (0, 4] \rightarrow \mathbb{R}$ be completely monotonic and $\mathcal{C} \in S^{m-1}$ be a sharp configuration. Then for all $\mathcal{C}' \in S^{m-1}$ with $|\mathcal{C}| = |\mathcal{C}'|$ we have $E_f(\mathcal{C}) \leq E_f(\mathcal{C}')$.

In this case we say that $\mathcal{C}$ is universally optimal. $\mathcal{C}$ minimises potential energy among all point configurations with $|\mathcal{C}|$ elements and for all completely monotonic potential function.

Applies to:

n	N	M	Inner products	Name
2	N	$N-1$	$\cos(2\pi j/N) \ (1 \leq j \leq N/2)$	N -gon
n	$N \leq n$	1	$-1/(N-1)$	simplex
n	$n+1$	2	$-1/n$	simplex
n	$2n$	3	$-1, 0$	cross polytope
3	12	5	$-1, \pm 1/\sqrt{5}$	icosahedron
4	120	11	$-1, \pm 1/2, 0, (\pm 1 \pm \sqrt{5})/4$	600-cell
8	240	7	$-1, \pm 1/2, 0$	E_8 roots
7	56	5	$-1, \pm 1/3$	kissing
6	27	4	$-1/2, 1/4$	kissing/Schläfli
5	16	3	$-3/5, 1/5$	kissing
24	196560	11	$-1, \pm 1/2, \pm 1/4, 0$	Leech lattice
23	4600	7	$-1, \pm 1/3, 0$	kissing
22	891	5	$-1/2, -1/8, 1/4$	kissing
23	552	5	$-1, \pm 1/5$	equiangular lines
22	275	4	$-1/4, 1/6$	kissing
21	162	3	$-2/7, 1/7$	kissing
22	100	3	$-4/11, 1/11$	Higman-Sims
q^2+1	$(q+1)(q^3+1)$	3	$-1/q, 1/q^2$	isotropic subspaces
	(4 if $q=2$)			(q a prime power)

See

Wu & Kumar

It does not apply to the 600-cell but very similar techniques can be used.

More conjectured universal optima:

n	N	t	References
10	40	$1/6$	Conway, Sloane, and Smith [Sl], Hovinga [H]
14	64	$1/7$	Nordstrom and Robinson [NR], de Caen and van Dam [dCvD], Ericson and Zinoviev [EZ]

see Ballinger et al

Question Are there all universal optima for S^{m-1} ?

Ingredients for proof of main theorem

- A) Fourier analysis on the sphere
(spherical harmonics)
- B) Spherical design strength test
- C) Linear programming bound
- D) Hermite interpolation
- E) Orthogonal polynomials

A) Spherical harmonics

With S^{m-1} one associates the family of orthogonal polynomials $P_k^n(t)$, $k=0,1,\dots$, where P_k^n is a univariate polynomial of degree k normalised by $P_k^n(1) = 1$ and satisfying the orthogonality relation

$$(xx) \quad \int_{-1}^1 P_k^n(t) P_l^n(t) (1-t^2)^{\frac{m-3}{2}} dt = 0 \quad \text{if } k \neq l$$

(see p. 99 of Co 2018/19).

→ We are going to explain the P_k^n in some detail → spherical harmonics.

→ Co 2018/19: Schoenberg's theorem.

About polynomials satisfying (xx)

If $n=2$, then P_k^2 are the Chebyshev polynomials

T_k with $T_k(\cos \theta) = \cos(k\theta)$

For larger n , P_k^n can be determined by Gram-Schmidt-orthogonalisation using the inner product

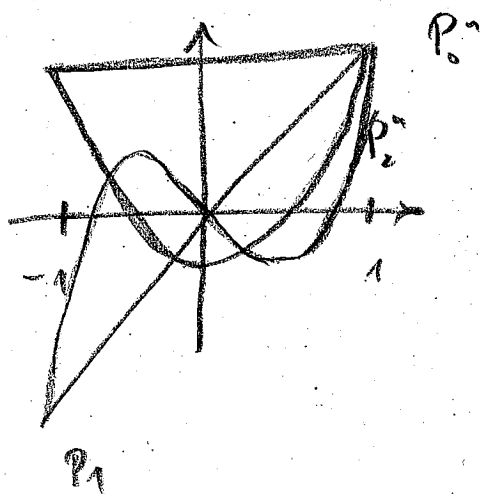
$$(f, g) = \int_{-1}^1 f(t) g(t) (1-t^2)^{\frac{n-3}{2}} dt.$$

Values for small k :

$$P_0^n(t) = 1$$

$$P_1^n(t) = t$$

$$P_2^n(t) = \frac{n}{n-1} t^2 - \frac{1}{n-1}$$



P_k^n are known by many computer packages:

Jacobi polynomials / Gegenbauer polynomials

(Warning! Be careful with parameter / normalisation)

Important property of P_k^n

Positivity property

the map $(x, y) \mapsto P_k^n(x \cdot y)$ for $x, y \in S^{n-1}$ is a positive kernel, i.e. the infinite matrix

$[P_k^n(x \cdot y)]_{x, y \in S^{n-1}}$ is positive semidefinite.

In other words,

$\forall N \in \mathbb{N} \forall x_1, \dots, x_N \in S^{n-1}$ the finite submatrix

$[P_k^n(x_i \cdot x_j)]_{i, j=1, \dots, N}$ is pos. semidef.

Proof: $\leadsto$ spherical harmonics.

b) Spherical design strength test

Lemma $\mathcal{C} \subseteq S^{m-1}$ spherical M -design

$$\Leftrightarrow \sum_{x, y \in \mathcal{C}} P_k^n(x \cdot y) = 0 \quad \text{for } k=1, \dots, M.$$

Proof $\leadsto$ spherical harmonics

d) Linear programming bound

for lower bounding f -potential energy

Theorem (Yudin, 1992)

Let $f: (0, 4]$ be any function. Suppose

$h: [-1, 1] \rightarrow \mathbb{R}$ is a polynomial such that

(i) $h(t) \leq f(2-2t)$ for all $t \in [-1, 1]$,

(ii) $h(t) = \sum_{k=0}^d \alpha_k P_k^n(t)$ with $\alpha_0, \dots, \alpha_d \geq 0$.

Then for every $\mathcal{C} \subseteq S^{m-1}$ with $|\mathcal{C}| = N$ we have

$$E_f(\mathcal{C}) \geq N^2 \alpha_0 - N h(1).$$

Remarks

- Yudin's bound generalises the LP method of Delsarte, Goethals, Seidel (1977) (see CO 2018/19: page 100) for spherical cap packings.
- Finding an optimal h , giving the best possible lower bound, is an (infinite-dimensional) LP.
- Yudin's bound generalises to unequally charged particles. $\leadsto$ Exercise.
- Break line of proof of main theorem: If $\mathcal{C}$ is a sharp configuration, then the LP bound is sharp (lower bound = upper bound) for all completely monotonic functions.

Proof (of LP bound)

Observe $\|x-y\|^2 = x \cdot x - 2x \cdot y + y \cdot y$
 $= 2 - 2x \cdot y$ for $x, y \in S^{n-1}$.

By assumption (i) and observation

$$E_f(\mathcal{C}) = \sum_{\substack{x, y \in \mathcal{C} \\ x \neq y}} f(\|x-y\|^2) \geq \sum_{\substack{x, y \in \mathcal{C} \\ x \neq y}} h(x \cdot y)$$

On the other hand, by assumption (ii) and the positivity property

$$\begin{aligned} \sum_{\substack{x, y \in \mathcal{C} \\ x \neq y}} h(x \cdot y) &= -N h(1) + \sum_{x, y \in \mathcal{C}} h(x \cdot y) \\ &= -N h(1) + \sum_{x, y \in \mathcal{C}} \sum_{k=0}^d \alpha_k P_k^*(x \cdot y) \\ &= -N h(1) + \sum_{k=0}^d \alpha_k \underbrace{\sum_{x, y \in \mathcal{C}} P_k^*(x \cdot y)}_{\geq 0} \\ &\geq -N h(1) + \alpha_0 \underbrace{\sum_{x, y \in \mathcal{C}} P_0^*(x \cdot y)}_{= N^2} = -N h(1) + \alpha_0 N^2. \end{aligned}$$



Remark If h proves sharp bound for $E_f(e)$, then

• $h(x \cdot y) = f(2 - 2x \cdot y)$ for all $x, y \in \mathcal{C}, x+y$.

• if $d_k > 0$, then $\sum_{x, y \in \mathcal{C}} P_k^n(x \cdot y) = 0$

($\rightarrow$ spherical design strength test).

D) Hermite interpolation

Idea Given a function f , find a polynomial (of low degree) that matches the values of f and some of its derivatives at some specified points.

Theorem Given $f \in \mathcal{C}^\infty([a, b])$, distinct points $t_1, \dots, t_m \in [a, b]$, positive integers $k_1, \dots, k_m$. Then there is a unique polynomial of degree less than

$$D = k_1 + k_2 + \dots + k_m \text{ such that}$$

$$p^{(k)}(t_i) = f^{(k)}(t_i) \quad \text{for all } 1 \leq i \leq m, \quad 0 \leq k < k_i.$$

Proof see your favorite numerics class / book.

Usage in the proof of main theorem

Interpolate f at position x, y for $x, y \in \mathbb{C}$ upto order 2 (i.e. value and 1st derivative) gives polynomial $h \rightsquigarrow$ assumption (i) fulfilled.

Orthogonal polynomials

- classical and useful topic in analysis (but not very fashionable...)
- very good books: for example
 - Szegő - Orthogonal polynomials, 1958
 - Andrews, Askey, Roy - Special functions, 1999
 - ...
- important properties: 3-term recurrence relation, interlacing of roots, ...

Let $w \in \mathcal{C}([0, b])$ be a continuous, nonnegative (weight) function, $w \neq 0$.

Def. A sequence of polynomials $(p_n)_{n=0,1,\dots}$, where $\deg p_n = n$, is orthogonal with respect to w if the following orthogonality relation holds:

$$\int_a^b p_m(x) p_n(x) w(x) dx = \delta_{mn}.$$

Theorem A sequence of orthogonal polynomials $(p_n)_n$ satisfies the 3-term recurrence relation

$$p_{n+1}(x) = (A_n x + B_n) p_n(x) + C_n p_{n-1}(x), \quad n \geq 0,$$

with $p_{-1}(x) = 0$ and $A_n, B_n, C_n \in \mathbb{R}$.

Proof Determine A_n so that

$$p_{n+1}(x) - A_n x p_n(x)$$

has degree n .

Then

$$p_{n+1}(x) - A_n x p_n(x) = \sum_{k=0}^n \beta_k p_k(x). \quad (*)$$

for $\beta_k \in \mathbb{R}$. If q is polynomial of degree $m < n$,
then by orthogonality

$$\int_a^b p_n(x) q(x) w(x) dx = 0.$$

Hence, $\beta_k = 0$ for $k < n-1$, because

$$\begin{aligned} & \int_a^b (p_{n+1}(x) - A_n x p_n(x)) p_k(x) w(x) dx \\ &= \int_a^b p_{n+1}(x) p_k(x) w(x) dx - A_n \int_a^b p_n(x) (x p_k(x)) w(x) dx \\ &= 0. \end{aligned}$$

So,

$$\begin{aligned} p_{n+1}(x) &= A_n x p_n(x) + \beta_n p_n(x) + \beta_{n-1} p_{n-1}(x) \\ &= \underbrace{(A_n x + \beta_n)}_{B_n} p_n(x) + \underbrace{\beta_{n-1}}_{C_n} p_{n-1}(x) \end{aligned}$$



Remark

- More can be said about A_n, B_n, C_n .
- 3-term recurrence relations characterise orthogonal polynomials ($\leadsto$ Favard's theorem).

Example

Chebyshev polynomials of the first kind

$$T_0(x) = 1$$

$$T_1(x) = x$$

$$T_{n+1}(x) = 2x T_n(x) - T_{n-1}(x).$$

$$T_n \text{ satisfies } T_n(\cos \theta) = \cos n\theta.$$

More about: D) Hermite Interpolation

Lemma (Remainder formula)

Given $f \in C^\infty([a, b])$, distinct points $t_1, \dots, t_m \in [a, b]$, positive integers $k_1, \dots, k_m$. Let p the corresponding Hermite interpolation polynomial of f , that is

$\deg p < D = k_1 + \dots + k_m$ and $p^{(k_i)}(t_i) = f^{(k_i)}(t_i)$

for all $1 \leq i \leq m$, $0 \leq k < k_i$. For every $t \in [a, b]$

there exists $z \in (a, b)$ with

$$\min\{t, t_1, \dots, t_m\} < z < \max\{t, t_1, \dots, t_m\}$$

so that

$$f(t) - p(t) = \frac{f^{(D)}(z)}{D!} (t-t_1)^{k_1} \dots (t-t_m)^{k_m}.$$

Proof Define

$$g(t) = \frac{f(t) - p(t)}{(t-t_1)^{k_1} \dots (t-t_m)^{k_m}}.$$

For $t \in [a, b] \setminus \{t_1, \dots, t_m\}$ consider the function

$$h(s) = f(s) - p(s) - g(t) (s-t_1)^{k_1} \dots (s-t_m)^{k_m}$$

We have $h^{(k_i)}(t_i) = 0$ for $1 \leq i \leq m, 0 \leq k < k_i$

and $h(t) = 0$. So h has $D+1$ roots (counted with multiplicity) in the interval

$$[\min\{t, t_1, \dots, t_m\}, \max\{t, t_1, \dots, t_m\}].$$

By iterated use of Rolle's theorem ($\rightarrow$ special case of mean value thm), there exists τ in this interval so that $h^{(D)}(\tau) = 0$. Hence,

$$h^{(D)}(\tau) = f^{(D)}(\tau) - g(t) D! = 0,$$

and the lemma follows. $\square$

Def $I \subseteq \mathbb{R}$ interval, $f \in C^\infty(I)$ is called absolutely monotonic if $f^{(k)}(x) \geq 0$ for all $x \in I, k \in \mathbb{N}$.

[strictly absolutely monotonic if $f^{(k)}(x) > 0$ for all $x \in \text{interior of } I, k \in \mathbb{N}$]

Systematic notation for Hermite interpolation

Let $f \in C^\infty$ and let g be a polynomial of $\deg g \geq 1$. By $H(f, g)$ denote the polynomial p with $\deg p < \deg g$ such that p agrees with f at each root of g up to the order of that root, i.e.

$$g^{(k_i)}(t_i) = 0 \Rightarrow p^{(k_i)}(t_i) = f^{(k_i)}(t_i).$$

Define

$$Q(f, g)(t) = \frac{f(t) - H(f, g)(t)}{g(t)}.$$

Q extends to a C^∞ -function at the roots of g .

Lemma Let f be absolutely monotonic in (a, b) , let $g(t) = (t-t_1)^{k_1} \dots (t-t_m)^{k_m}$ with $t_1, \dots, t_m \in [a, b]$.

Then $Q(f, g)(t)$ is absolutely monotonic. The same is true for strictly absolutely monotonic.

Proof $\rightarrow$ Exercise (see Proposition 2.2 of the paper)

More about E) Orthogonal polynomials

Def Let $w \in \mathcal{C}([a, b])$ be a continuous weight function. We say w is positive definite up to degree N if for all polynomials f with $\deg f \leq N$ we have

$$\int_a^b f(t)^2 w(t) dt \geq 0,$$

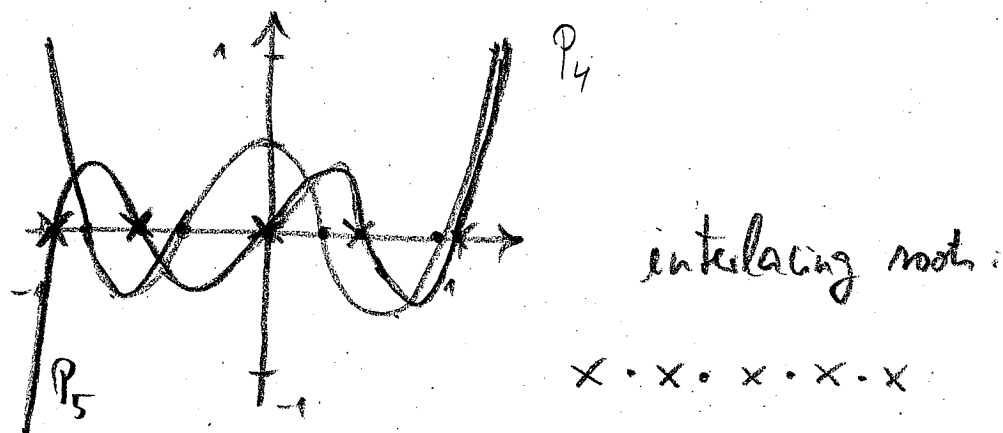
with equality $\Leftrightarrow f = 0$.

Lemma If w is positive definite up to degree N , then there are unique monic poly. $q_0, \dots, q_{N+1}$ such that $\deg(q_i) = i$ and

$$\int_a^b q_i(t) q_j(t) w(t) dt = 0 \quad \text{for } i \neq j.$$

For each i , q_i has i distinct real roots. The roots of q_i and q_{i-1} are interlaced.

Example Legendre polynomials P_4, P_5



Proof (of lemma)

• As usual one can determine q_i by applying Gram-Schmidt orthogonalization of $1, t, t^2, \dots, t^{N+1}$:

$$q_i(t) = t^i - \sum_{j=0}^{i-1} q_j(t) \frac{\int_a^b t^i q_j(s) w(s) ds}{\int_a^b q_j(s)^2 w(s) ds},$$

because $\langle p, q \rangle = \int_a^b p(t) q(t) w(t) dt$ determines a nondegenerate inner product on the space of polynomials of degree $\leq N$.

Hence, the q_i 's are uniquely determined.

• They satisfy a 3-term recurrence relation:

The polynomial $q_i(t) - t q_{i-1}(t)$ has degree $i-1$ and

is orthogonal to all polynomials of degree $< i-2$.
Hence there are constants a_i and b_i with

$$q_i(t) = (t + a_i) q_{i-1}(t) + b_i q_{i-2}(t).$$

• We have $b_i < 0$ because

$$\begin{aligned} b_i \int_a^b q_{i-2}(t)^2 w(t) dt &= \int_a^b (q_i(t) q_{i-2}(t) - (t + a_i) q_{i-1}(t) q_{i-2}(t)) w(t) dt \\ &= \int_a^b -t q_{i-1}(t) q_{i-2}(t) w(t) dt \end{aligned}$$

and

$$\int_a^b q_{i-1}(t) t q_{i-2}(t) w(t) dt = \int_a^b q_{i-1}(t)^2 w(t) dt.$$

• Show by induction that the roots of q_i and q_{i-1} are real and interlaced.

Base case : $i = 1$: clear.

Induction step Suppose roots of q_{i-1} and q_{i-2} are real and interlaced.

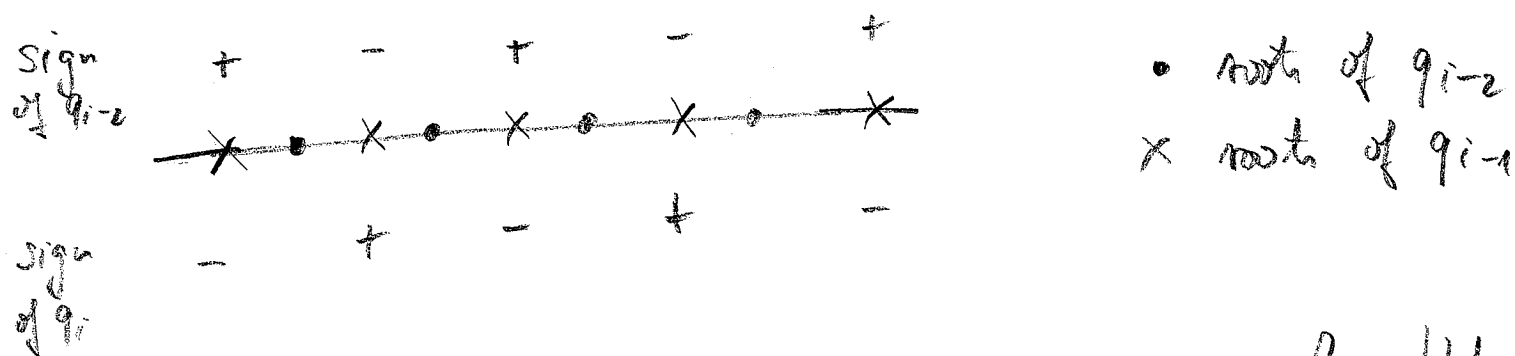
Let r be a root of q_{i-1} . By (3TRR)

$$q_i(r) = b_i q_{i-2}(r).$$

By induction: q_{i-2} alternates in sign at roots of q_{i-1} .

Since $b_i < 0$, q_i and q_{i-2} have opposite signs at roots of q_{i-1} . Hence, q_i has a root between each pair of consecutive roots of q_{i-1} .

Picture of the situation:



Because q_i and q_{i-2} have the same sign, when $|t|$ is sufficiently large, q_i must have a root greater than every root of q_{i-1} and a root less than every root of q_{i-1} . This finishes the proof since all i roots of q_i are at the promised position.

Lemma (Gauss-Jacobi quadrature)

Let $w \in C([a, b])$ be positive definite up to every degree $N \in \mathbb{N}$, which defines the monic, orthogonal polynomials $(p_n)_{n \in \mathbb{N}}$. Let $\alpha \in \mathbb{R}$ and let

$$\tau_1 < \tau_2 < \dots < \tau_n$$

be the roots of $p_n + \alpha p_{n-1}$. Then there exist $\lambda_1, \dots, \lambda_n > 0$ so that for every polynomial f , $\deg f \leq 2n-2$:

$$\int_a^b f(t) w(t) dt = \sum_{i=1}^n \lambda_i f(\tau_i). \quad (67Q)$$

Proof There exist $\lambda_1, \dots, \lambda_n \in \mathbb{R}$ with

$$\int_a^b f(t) w(t) dt = \sum_{i=1}^n \lambda_i f(\tau_i) \quad (*)$$

for all polynomials f , $\deg f < n$.

Became Apply Lagrange interpolation with nodes $\pi_1, \dots, \pi_n$. This gives

$$f(t) = \sum_{i=1}^n l_i(t) f(\pi_i) \quad \text{with Lagrange polynomials } l_i$$

($l_i(\pi_j) = \delta_{ij}$).

Then

$$\int f(t) w(t) dt = \sum f(\pi_i) \underbrace{\int l_i(t) w(t) dt}_{= \lambda_i}.$$

Let f be a polynomial with $\deg f \leq 2n-2$.

Use polynomial division to find polynomials g, h so that

$$f = (p_n + \alpha p_{n-1}) g + h$$

with $\deg h < n$ and $\deg g \leq n-2$. Then by the orthogonality relation

$$\int f(t) w(t) dt = \int h(t) w(t) dt$$

and furthermore $f(\pi_i) = h(\pi_i)$.

So (*) extends to all polynomials f with $\deg f \leq 2n-2$.

To show: $\lambda_i > 0$

For fixed i specialize f to $f(t) = \prod_{j \neq i} (t - \tau_j)^2$

Then $\deg f = 2n-2$ and by (*)

$$\underbrace{\int_0^1 f(t) v(t) dt}_{>0} = \lambda_i f(\tau_i) = \lambda_i \underbrace{\prod_{j \neq i} (\tau_i - \tau_j)^2}_{>0}$$

$$\Rightarrow \lambda_i > 0.$$

□

Again: $w \in \mathcal{C}([a, b])$ positive definite up to
 every degree $N \in \mathbb{N}$. Defines $(p_n)_{n \in \mathbb{N}}$ orth. poly.

Let $\tau_1 < \dots < \tau_n$ be the roots of $p_n + \alpha p_{n-1}$.

For $j \in \{0, \dots, n\}$ define $w_j \in \mathcal{C}([a, b])$ by

$$w_j(t) = \prod_{i=0}^{j-1} (\tau_{n-i} - t) w(t).$$

Clearly, $w_0 = w$.

Lemma For $j \in \{0, \dots, n-1\}$, the weight function w_j
 is positive definite up to degree $n-j-1$.

Proof By (6.7.2)

$$\begin{aligned} \int f(t)^2 w_j(t) dt &= \int f(t)^2 \prod_{i=0}^{j-1} (\tau_{n-i} - t) w(t) dt \\ &= \sum_{i=1}^n \lambda_i \underbrace{f(\tau_i)}_{\geq 0} \underbrace{\prod_{l=0}^{j-1} (\tau_{n-l} - \tau_i)}_{\geq 0} \geq 0 \end{aligned}$$

for all polynomials f with

$$2 \cdot \deg f + j \leq 2n - 2 \Rightarrow \deg f \leq n - 1 - j/2.$$

We have equality " $= 0$ " iff

$$f(\tau_1) = f(\tau_2) = \dots = f(\tau_{n-j}) = 0.$$

If $\deg f \leq n - j - 1$, then $f = \bar{0}$. □

* Theorem For $k < n$, the polynomial

$\prod_{i=1}^k (t - \tau_i)$ can be expressed as

$$\prod_{i=1}^k (t - \tau_i) = \sum_{n=0}^k \alpha_n p_n \quad \text{with } \alpha_0, \dots, \alpha_k \geq 0.$$

Proof Consider weight function

$$w_j(t) = \prod_{i=0}^{j-1} (\tau_{n-i} - t) w(t), \quad 0 \leq j \leq n$$

Let $q_{j,i}$ be the monic orthogonal polynomial of deg i for w_j . ($q_{0,i} = p_i$, p_i orth. for w).

We have

$$q_{j,n-j}(t) = (t - \tau_1) \dots (t - \tau_{n-j})$$

because for a polynomial q with $\deg q < n-j$

$$\int_a^b q(t) q_{j,n-j}(t) w_j(t) dt$$

$$= \int_a^b q(t) (t - \tau_1) \dots (t - \tau_{n-j}) (t - \tau_{n-j+1}) \dots (t - \tau_n) w(t) dt =$$

$$w(t) dt =$$

$$= \int_a^b q(t) (p_n + \alpha p_{n-1})(t) w(t) dt$$

$$= 0.$$

So, for $i < n-j$, the largest root of $q_{j,i}$ is less than τ_{n-j} (by interlacing).

For $i \leq n-j$, the largest root of $q_{j-1,i}$ is less than τ_{n-j+1} , so $q_{j-1,i}(\tau_{n-j+1}) \neq 0$.

Hence for $i \leq n-j$ there are constants $\alpha_{j,i}$ so that

$$q_{j,i}(t) = \frac{q_{j-1,i+1}(t) + \alpha_{j,i} q_{j-1,i}(t)}{t - \tau_{n-j+1}}$$

because for $\alpha_{j,i}$ chosen so that

$$q_{j-1,i+1}(\tau_{n-j+1}) + \alpha_{j,i} q_{j-1,i}(\tau_{n-j+1}) = 0$$

we have for every polynomial q with $\deg(q) \leq i-1$

$$\begin{aligned}
& \int_a^b q \frac{q_{j-1,i+1}(t) + \alpha_{j,i} q_{j-1,i}(t)}{t - \tau_{n-j+1}} w_j(t) dt \\
&= \int_a^b q (q_{j-1,i+1}(t) + \alpha_{j,i} q_{j-1,i}(t)) w_{j-1}(t) dt \\
&= 0
\end{aligned}$$

Now the theorem follows from the next lemma applied to $q_{k,n-k}$.

Lemma (Lemma 3.8 in Cohn-Kumar paper)

For $1 \leq j \leq n$ and $i \leq n-j$ the polynomial $q_{j,i}$ is a positive linear combination of $q_{j-1,0}, \dots, q_{j-1,i}$.

Proof. $\rightarrow$ Exercise.



Choosing h for the LP bound

Let $f: (0, 4] \rightarrow \mathbb{R}$ be completely monotonic.

Define $a(t) = f(2-2t)$ (because $\|x-y\|^2 = 2-2x \cdot y$)

Function a is absolutely monotonic on $[-1, 1)$.

Three inputs to construct h :

$$n \in \mathbb{N} \quad (\text{for } S^{n-1})$$

$$m \in \mathbb{N}, \quad m \geq 1$$

$$\alpha \in \mathbb{R}, \quad \alpha \geq 0.$$

Let $t_1 < \dots < t_m$ be the roots of

$$P_m^m + \alpha P_{m-1}^m. \quad \text{We require } t_1 \geq -1 \text{ and then}$$

$$\{t_1, \dots, t_m\} \subseteq [-1, +1). \quad (\text{Szegő: } \alpha \leq \frac{-P_m^m(-1)}{P_{m-1}^m(-1)})$$

Let h be the Hermite interpolating polynomial of a of order 2 at each t_i .

$$h(t_i) = a(t_i) \quad \text{and} \quad h'(t_i) = a'(t_i) \quad \text{for } 1 \leq i \leq m.$$

Lemma $h(t) \leq a(t)$ for all $t \in [-1, 1)$.

Proof By the remainder formula (page 19)
there is z with

$$\min \{t, t_1, \dots, t_m\} < z < \max \{t, t_1, \dots, t_m\}$$

and

$$a(t) - h(t) = \underbrace{\frac{a^{(2m)}(z)}{(2m)!} (t-t_1)^2 \cdots (t-t_m)^2}_{\geq 0}$$

$$\Rightarrow a(t) \geq h(t).$$

□

Define

$$F(t) = P_m^n(t) + d P_{m-1}^n(t) = \prod_{i=1}^m (t-t_i)$$

Then, $h = H(a, F^2)$.

Def. A nonconstant polynomial with all its roots in $[-1, 1)$ is called conductive if for every absolutely monotonic function a on $[-1, 1)$

$H(a, g)$ lies in the convex cone spanned by P_m^k ,
 $m = 0, 1, \dots$ ($H(a, g)$ positive definite)

g is strictly conductive if it is conductive and for
all strictly absolutely monotonic a ,

$$H(a, g) = \sum_{m=0}^{\deg g - 1} \alpha_m P_m^k \text{ with } \alpha_m > 0.$$

($H(a, g)$ positive definite of degree
 $\deg g - 1$)

Lemma (a) g_1, g_2 conductive, g_1 positive definite, then
 $g_1 g_2$ conductive.

(b) g_1 conductive and strictly positive definite and
 g_2 strictly conductive, then $g_1 g_2$ strictly conductive.

Proof (of a)

$$H(a, g_1 g_2) = \underbrace{H(a, g_1)}_{\text{pos. def.}} + \underbrace{g_1}_{\text{pos. def.}} \underbrace{H(Q(a, g_1), g_2)}_{\substack{\text{absolutely} \\ \text{monotonic} \text{ (Prop. 2.1)}}}$$

(see proof of Proposition 2.2)

Known: The cone of positive definite functions is closed under taking (pointwise) product.

($\leadsto$ Schur-Hadamard product for matrices).

Proof of (1): similar. □

Lemma $H(a, F^2)$ is conductive

Proof. For $r \in [-1, 1)$, let

$$l_r(t) = t - r.$$

We have $H(a, l_r) = a(r) \geq 0$, so l_r is conductive (and strictly conductive if $r \neq -1$).

By the previous lemma: g conductive and positive definite
 $\Rightarrow g l_r$ conductive.

Consider $\prod_{i=1}^j (t - t_i)$ for $j \leq m$.

By the * Theorem (p. 31) the product for $j < m$ is positive semidefinite (even strictly). For $j = m$ the product is positive definite because $\alpha \geq 0$. Hence, by the previous lemma, each product is conductive. Hence, F^2 is conductive by the previous lemma. $\square$

Corollary $h = H(a, F^2)$ is positive definite.

Proof of main theorem

Let $f : (0, 4] \rightarrow \mathbb{R}$ be completely monotonic,

$a(t) = f(2 - 2t)$ absolutely monotonic,

let $\mathcal{C} \subseteq S^{n-1}$ be a sharp configuration with

$N = |\mathcal{C}|$. Let

$$-1 \leq t_1 < \dots < t_m < 1$$

be the inner products between distinct points in $\mathcal{C}$.

Let $h(t)$ be the Hermite interpolating polynomial that agrees with $a(t)$ to order 2.

• We already proved that $h(t) \leq a(t)$ for $t \in [-1, 1)$ (Lemma on page 35).

• Write $h(t) = \sum_{k=0}^d \alpha_k P_k^n(t)$, with $d = \deg h$.

We have $d = 2m - 1$. Now all inequalities in the proof of the linear programming bound

are equalities in this case since $h(t_i) = a(t_i)$
and since

$$\sum_{x, y \in \mathcal{C}} P_k^n(x, y) = 0 \quad \text{for all } 0 < k \leq 2m-1$$

because $\mathcal{C}$ is a spherical $(2m-1)$ -design.

$$\text{Hence, } E_f(\mathcal{C}) = N^2 \alpha_0 - N h(1)$$

under the condition $\alpha_0, \dots, \alpha_{2m-1} \geq 0$.

• h is positive definite.

Define $F(t) = \prod_{i=1}^m (t - t_i)$. Then $h = H(a, F^2)$.

Claim F^2 is conductive.

Lemma F is strictly positive definite.

Proof Set $F(t) = \sum_{k=0}^m \beta_k P_k^n(t)$. Then $\beta_m > 0$.

For $0 \leq k < m$ consider

$$\int_{S^{m-1}} F(x \cdot y) P_k^n(x \cdot y) d\omega(x) = \underbrace{C_k}_{>0} \beta_k.$$

for $y \in S^{m-1}$. Choose $y \in \mathcal{C}$. We have

$\deg(F) + k \leq 2 \deg(F) - 1$. Because $\mathcal{C}$ is a spherical $(2 \deg(F) - 1)$ -design, it follows that the integral is a positive constant times

$$\sum_{x \in \mathcal{C}} F(x \cdot y) P_k^n(x \cdot y).$$

By the definition of $\mathcal{C}$:

$$\sum_{x \in \mathcal{C}} F(x \cdot y) P_k^n(x \cdot y) = F(1) P_k^n(1) > 0.$$

□

Let $p_0, p_1, \dots$ be the monic orthogonal polynomials for the weight function $w(t) = (1-t)(1-t^2)^{\frac{n-3}{2}}$.

One can show that these polynomials lie in the convex cone generated by $P_0^n, P_1^n, \dots$

Lemma $\exists \alpha : F = p_m + \alpha p_{m-1}$

Proof To show

$$(*) \int_{-1}^1 (1-t) F(t) q(t) (1-t^2)^{\frac{n-3}{2}} dt = 0$$

for all polynomials q with $\deg q \leq m-2$.

Because $\mathcal{C}$ is a spherical $(2m-1)$ -design,
the integral $(*)$ equals a positive linear combination
of

$$(1-t_1) F(t_1) p(t_1), \dots, (1-t_m) F(t_m) p(t_m),$$

$$\text{and } (1-1) F(1) p(1).$$

Thus, $(*) = 0$.

□

This lemma combined with the Theorem on page 31

shows that $\prod_{i=1}^j (t-t_i)$ for $j < m$

is positive definite. The case $j=m$ is the fact

that F is positive definite.

Now the proof that F and F^2 are conjugative follows as in the previous paragraph on choosing h for the LP bound. □

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