

Convex Optimisation (WS 2015/16)

Chapter I Introduction

mathematical program ($\hat{=}$ mathematical optimisation problem)

given: functions $f_0, f_1, \dots, f_m : \mathbb{R}^n \rightarrow \mathbb{R} \cup \{\infty\}$,
constants $b_1, \dots, b_m \in \mathbb{R}$

find: minimise $f_0(x)$
such that $x \in \mathbb{R}^n$
 $f_j(x) \leq b_j, j = 1, \dots, m.$

Vocabulary: f_0 : objective function; f_1, \dots, f_m : constraint

functions; x : optimisation variable; $x \in \mathbb{R}^n$ satisfying

$f_j(x) \leq b_j$ for all $j = 1, \dots, m$: feasible solution; feasible
solution x with $f_0(x) \leq f_0(x')$ for all feasible solution x' :

optimal solution

Important examples: Linear programs (LPs)

- all functions f_0, \dots, f_m are linear.

- LPs are well-understood in theory and practice
(\leadsto OR)

Nonlinear programs (NLPs)

One of the functions f_0, \dots, f_m is not linear:

$$f_0(x, y) = e^{\sin(50x)} + \sin(60e^y) + \sin(70 \sin x) \\ + \sin(\sin(80y)) - \sin(10(x+y)) + \frac{x^2 + y^2}{4}$$

optimal solution $(-0, 0244\dots, 0, 2106\dots)$

However, NLPs are too general for developing a useful mathematical theory.

Convex programs

all functions f_0, \dots, f_m are convex.

Recall:

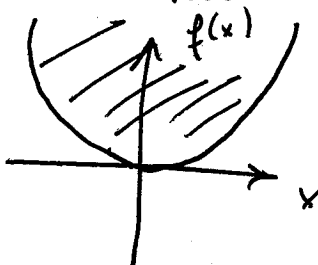
Def. 1 A function $f: \mathbb{R}^n \rightarrow \mathbb{R} \cup \{\infty\}$ is called convex

if the epigraph

$$\text{epi}(f) = \{ (x, \alpha) \in \mathbb{R}^{n+1} : f(x) \leq \alpha \}$$

is a convex set.

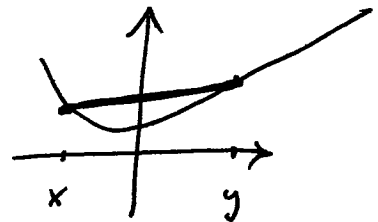
Ex. 2 $f(x) = x^2$ is convex

$\text{epi}(f) =$ 

Proposition 3 A function $f: \mathbb{R}^n \rightarrow \mathbb{R} \cup \{\infty\}$ is convex if and only if f satisfies Jensen's inequality:

$$f(tx + (1-t)y) \leq tf(x) + (1-t)f(y)$$

for all $t \in [0, 1]$, $x, y \in \mathbb{R}^n$.



Proof " \Rightarrow ": Define $\tilde{x} = (x, f(x))$, $\tilde{y} = (y, f(y)) \in \mathbb{R}^{n+1}$.

Clearly, $\tilde{x}, \tilde{y} \in \text{epi}(f)$. Hence, $t\tilde{x} + (1-t)\tilde{y} \in \text{epi}(f)$ for $t \in [0, 1]$. So,

$$(tx + (1-t)y, tf(x) + (1-t)f(y)) \in \text{epi}(f),$$

and

$$f(tx + (1-t)y) \leq tf(x) + (1-t)f(y).$$

" \Leftarrow ": very similar. \square

Useful convexity criterion (OR 2015, Ex. 5.4)

Proposition 4 Let $f: \mathbb{R}^n \rightarrow \mathbb{R}$ be a function which is twice continuously differentiable. Function f is convex if and only if its Hessian

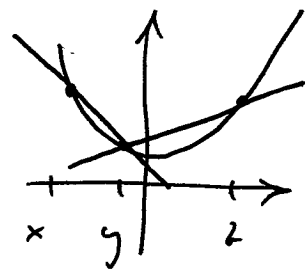
$$H(f)(z) = \left(\frac{\partial^2}{\partial x_i \partial x_j} f(z) \right)_{i,j=1,\dots,n} \in \mathbb{R}^{n \times n}$$

is positive semidefinite for all $z \in \mathbb{R}^n$.

[Schwarz's theorem implies that $H(f)$ is symmetric.]

Proof First consider the case $n=1$. A function $f: \mathbb{R} \rightarrow \mathbb{R}$ is convex if and only if for all $x, y, z \in \mathbb{R}$ with $x < y < z$ inequality

$$(*) \quad \frac{f(y) - f(x)}{y - x} \leq \frac{f(z) - f(y)}{z - y}$$



holds. By the mean value theorem there are $\xi_1 \in [x, y]$, $\xi_2 \in [y, z]$ such that

$$f'(\xi_1) = \frac{f(y) - f(x)}{y - x} \quad \text{and} \quad f'(\xi_2) = \frac{f(z) - f(y)}{z - y}$$

hold. Suppose f is convex, then this together with (*) implies

that f' is monotonically increasing. Hence, $f'' \geq 0$.

Suppose $f'' \geq 0$. Then f' is monotonically increasing and (*) is fulfilled. Hence, f is convex.

Now the multiple dimensional case: f is convex

iff for all $x, y \in \mathbb{R}^n$ the function $g_{xy}(\alpha) = f(\alpha x + (1-\alpha)y)$

is convex. To show: $g''_{xy} \geq 0 \iff H_f(z)$ is positive semi-definite.

Apply chain rule to see this:

$$g'_{xy}(\alpha) = \sum_{i=1}^n \frac{\partial f}{\partial x_i} (y + \alpha(x-y)) (x_i - y_i)$$

$$g''_{xy}(\alpha) = \sum_{i=1}^n \sum_{j=1}^n \frac{\partial^2 f}{\partial x_i \partial x_j} (y + \alpha(x-y)) (x_i - y_i) (x_j - y_j)$$

□

fundamental properties of convex programs

$$\min f_0(x)$$

$$x \in \mathbb{R}^n$$

$$f_j(x) \leq b_j, \quad (f_0, \dots, f_m \text{ convex}).$$

i) Set of feasible solutions is convex.

[immediately from Jensen's inequality].

ii) Local minima are global minima = optimal solution.

Proposition 5 Consider the convex program

$$\begin{aligned} \min f_0(x) \\ x \in \mathbb{R}^n \\ f_j(x) \leq b_j, \quad j=1, \dots, m, \end{aligned}$$

where $C = \{x \in \mathbb{R}^n : f_j(x) \leq b_j, j=1, \dots, m\}$ is the (convex) set of feasible solutions. Let x_0 be a local minimizer of the convex program, i.e. there is an $\varepsilon > 0$ s.t.

$$f_0(x_0) = \inf \{ f_0(y) : y \in C, \|x_0 - y\| \leq \varepsilon \}.$$

Then x_0 is an optimal solution of the convex program.

Proof Assume x_0 is not an optimal solution. Then there is a feasible solution $z \in C$ with $f(z) < f(x_0)$. Clearly, $\|x_0 - z\| > \varepsilon$. Define $y \in [x_0, z]$ by

$$y = (1-\alpha)x_0 + \alpha z \quad \text{with } \alpha = \frac{\varepsilon}{\|x_0 - z\|} \in [0, 1]$$

which lies in C . Then $\|x_0 - y\| = \varepsilon$ and by the convexity of f_0 : $f(y) = (1-\alpha)f(x_0) + \alpha f(z) < f(x_0)$. ∇